
Algorithm 1: function $Z(T, B)$, $T = T(G)$, for $G \in \text{sDH}(2)$

Input: $T(G)$ is assumed rooted, B is a component of $T(G)$ **Output:** the pair $[\zeta(B), \zeta_{\bar{B}}(m)]$

- 1 Let $\{B_1, B_2, \dots, B_k\}$ the children of B connected to B by edges (m_i, m'_i) ,
 $i = 1, \dots, k$, where $m_i \in B$
 - 2 Let m the marked vertex connecting B to its parent in T , if it exists.
 - 3 Let $a = b = c = d = e = f = \langle 0, 0 \rangle$
 - 4 $a = \max\{\zeta(u, v) \mid (u, v) \text{ is a stretch-pair of } B\}$
 - 5 $b = \max\{\text{first}(Z(T, B_i)) \mid i = 1, \dots, k\}$
 - 6 $c = \max\{\zeta_B(m_i) + \text{second}(Z(T, B_i)) \mid i = 1, \dots, k\}$
 - 7 $d = \max\{\zeta(m_i, m_j) + \text{second}(Z(T, B_i)) +$
 $\text{second}(Z(T, B_j)) \mid (m_i, m_j) \text{ is a stretch-pair of } B\}$
 - 8 **if** m *exists* **then**
 - 9 $e = \zeta_B(m)$
 - 10 $f = \max\{\zeta(m, m_i) + \text{second}(Z(T, B_i)) \mid (m, m_i) \text{ is a stretch-pair of } B\}$
 - 11 **return** $[\max\{a, b, c, d\}, \max\{e, f\}]$
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Algorithm 2: computing $s(G)$

Input: a graph G with any component B of $\mathcal{D}(G)$ such that $s(B) < 2$ **Output:** $s(G)$

- 1 compute $\mathcal{D}(G)$ and $T(G)$
 - 2 choose a component B of $\mathcal{D}(G)$ as a root of $T(G)$
 - 3 return $\sigma(\text{first}(Z(T(G), B)))$
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